

Finite Horizon Dynamic Programming: Getting Value from Spending Symmetry

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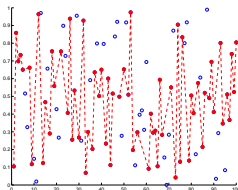
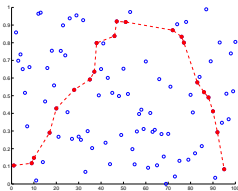
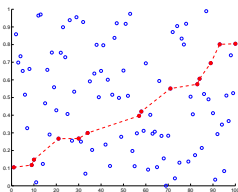
Stochastic Processes and Applications, Buenos Aires, August 8, 2014

Some History and Motivation

- Famous combinatorial problems with long mathematical history on sequences of n real numbers, or permutations of the integers $1, \dots, n$
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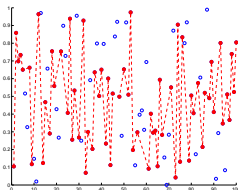
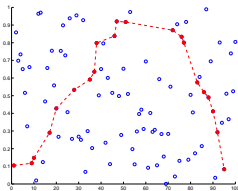
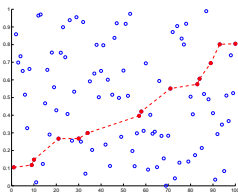
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 - ▶ **Longest** monotone subsequences: Hammersley (1972), Kingman (1973), Logan and Shepp (1977), Veršik and Kerov (1977), ...
 - ▶ **Longest** Unimodal subsequences: Steele (1981)
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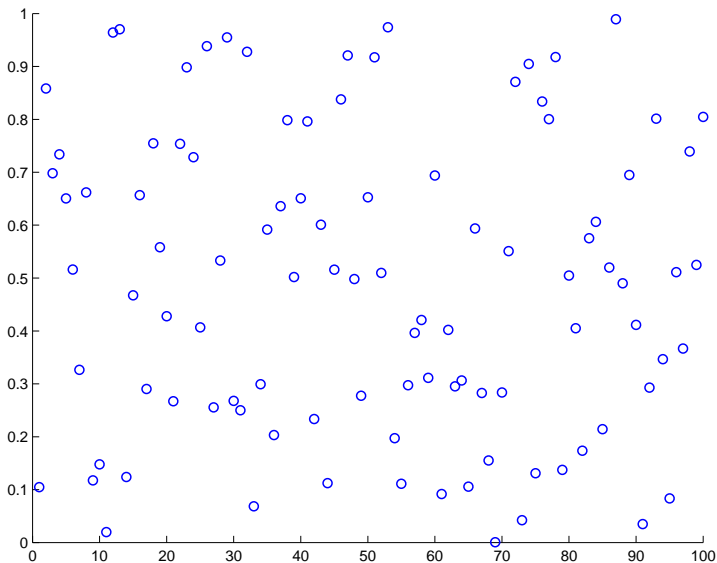


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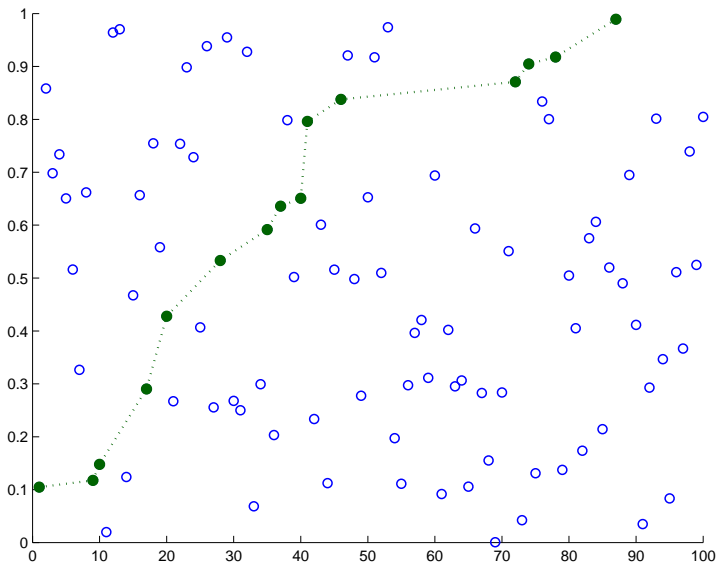
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- Now ... Study the sequential (**on-line**) version of these problems
 - ▶ **Objective:** maximize the expected length (number of selections) of monotone, unimodal and alternating subsequences



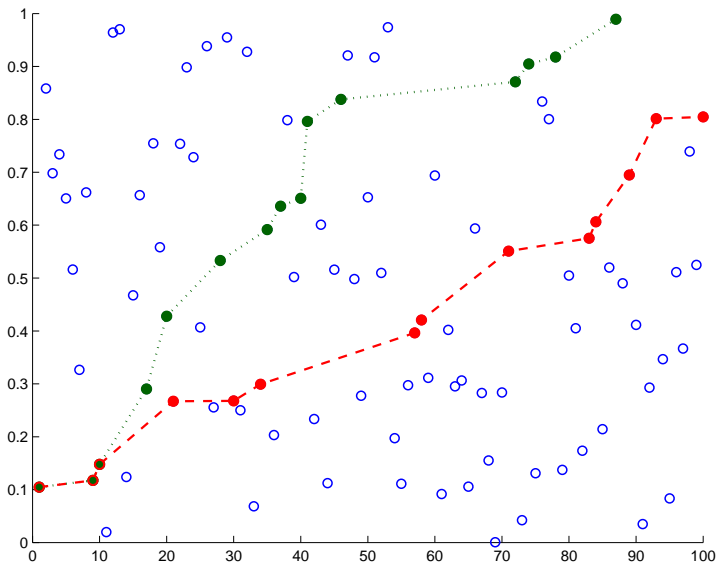
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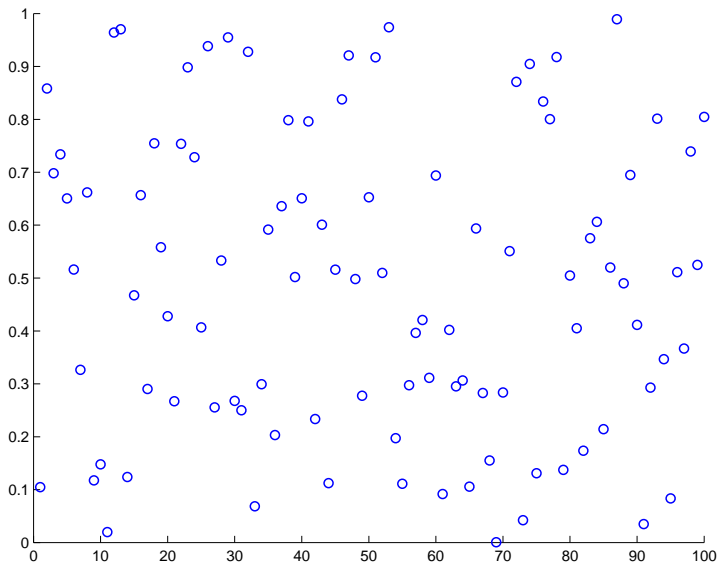
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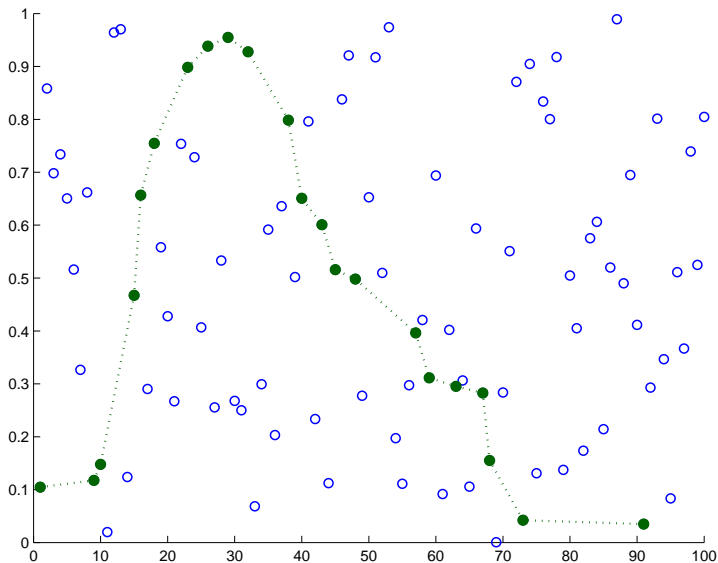
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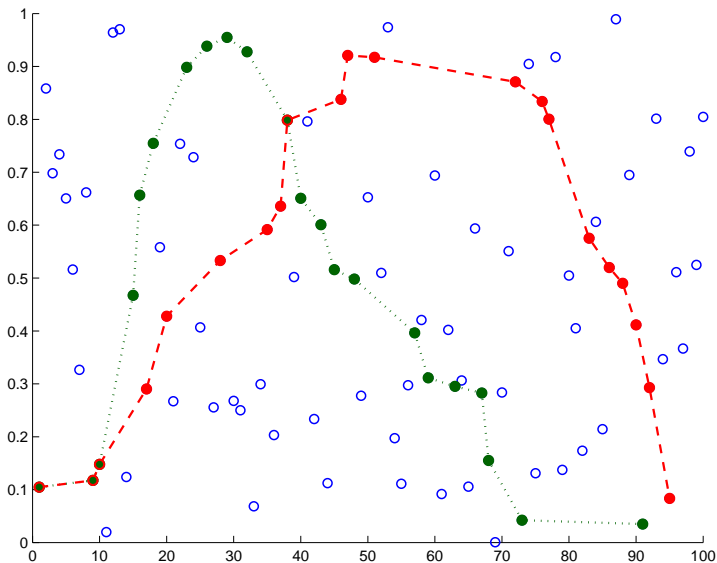
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$U_n^o(\pi_n^*) = 21$



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- There is a CLT for the On-Line Alternating Subsequence Problem (briefly noted in next frame)
- There has much further work on the *On-Line Selection of a Monotone Increasing Subsequence*, the original motivating problem. This will get most of our attention.

Sequentially Selected Alternating Series — A CLT

Theorem (Arlotto & Steele, AAP 2014)

There is a constant $\sigma > 0$ such that

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- **Source of Juice?** Very detailed analytical understanding of the acceptance threshold functions.

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There is a policy $\pi^ \in \Pi(n)$ such that*

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- **Puzzle:** A CLT is far from a sure thing. For the off-line problem one does NOT have a CLT — One has the famous Tracy-Widom Law.

Poissonization: A Homogenizing Trick with Benefits

- If one takes a sample size $N(t)$ that is Poisson with mean t there are several benefits: (a) optimal policies are stationary — no horizon effects and (b) one gets the machinery of infinitesimal generators, Dynkin Martingale, etc. There is long history of applications, perhaps starting with Lucien LeCam.

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 - ▶ Only “one of the five steps” to the proof of the CLT for the finite horizon LIS uses what one could call classical de-Poissonization.

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 - 5 The CLT itself

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Out of Five: Only One for Free

- The CLT of Bruss and Delbean has five parts:
 - ① Mean lower bound: $(2t)^{1/2} - O(\log(t))$
 - ② Mean upper bound: $(2t)^{1/2}$
 - ③ Variance lower bound: $\frac{1}{3}(2t)^{1/2} - O(1)$
 - ④ Variance upper bound: $\frac{1}{3}(2t)^{1/2} + O(\log t)$
 - ⑤ The CLT itself
- Only one of these steps has what one can properly call a de-Poissonization.
- De-Poissonization gives us the mean lower bound for the finite horizon problem — and leaves us four steps to go.

De-Poissonization of the Mean Lower Bound: One Proof

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- The Poisson strategy is a *suboptimal strategy* for a problem where one knows ex-ante that the sample has size j , so we have

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- But we have another property: the map $\phi(j) = \mathbb{E}[L_j^\circ]$ is concave. Jensen's inequality then forks up

$$\mathbb{E}[L_{N(n)}^\circ] \leq \sum_{j=0}^{\infty} e^{-n} \frac{n^j}{j!} \mathbb{E}[L_j^\circ] \leq \mathbb{E}[L_n^\circ].$$

Thus, we have *lossless transference* of any mean lower bound from the Poisson model to the Finite Horizon model.

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- How about the upper bound for $\text{Var}[L_n^o]$?
- Alessandro and I were stuck here for a long time.

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- The flood gate is opened and more analysis of the same flavor (but with plenty of details) lead us through the Martingale CLT to a CLT for the Finite Horizon Selection Problem for LIS:

$$\frac{3^{1/2}\{L_n^o - (2n)^{1/2}\}}{(2n)^{1/4}} \implies N(0, 1).$$

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● ¡Gracias por su atención !

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